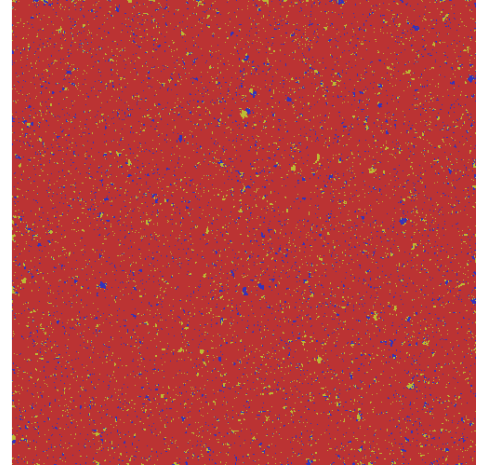
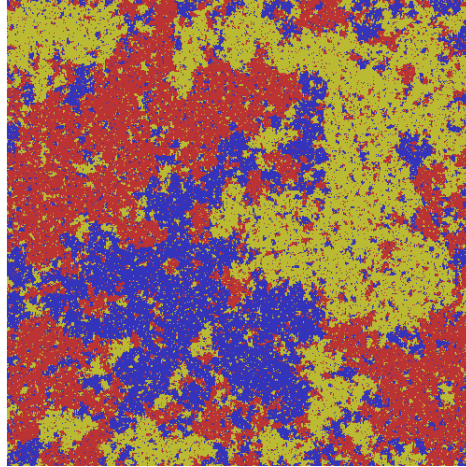
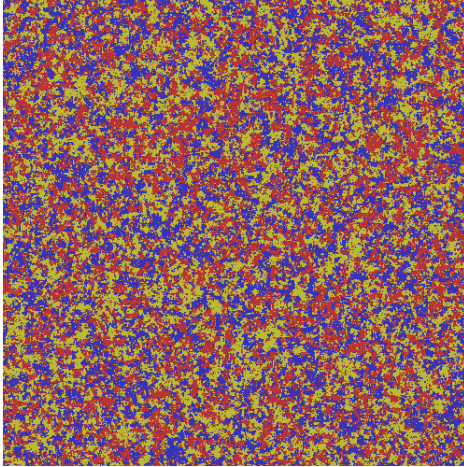


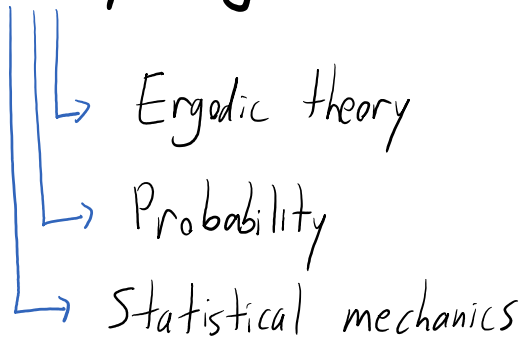
PHASE TRANSITIONS AND FINITARY CODINGS

Yinon Spinka

Phase transitions



Finitary codings



Akçoglu, Denker, Fiebig, del Junco, Kalikow, Keane, Krieger, Rahe, Rudolph, Schmidt, Serafin, Smorodinsky, Weiss

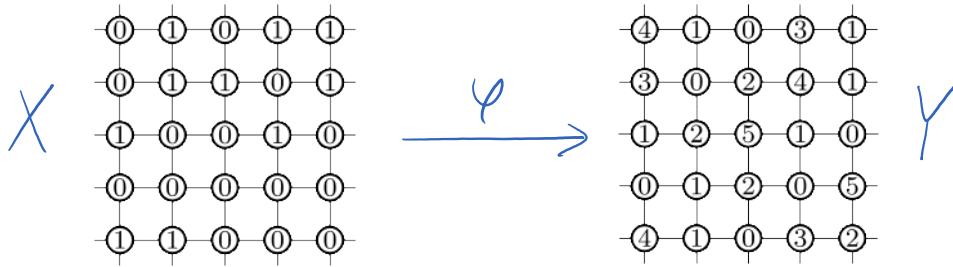
Angel, Benjamini, Gurel-Gurevich, Harvey, den Hollander, Holroyd, Hutchcroft, Keane, Kosloff, Levy, Lyons, Meyerovitch, Nazarov, Peled, Peres, Romik, Schramm, Soo, Steif, Timar, Wilson

van den Berg, Häggström, Harel, Ray, Steif

Closely related to → Phase transitions
→ Algorithms for perfect sampling (exact simulation)

Introduction:

- $X = (X_u)_{u \in \mathbb{Z}^d}$, $Y = (Y_u)_{u \in \mathbb{Z}^d}$ (translation invariant) [eg. X IID, Y trans-inv. Gibbs measure]



- Y is a **factor** of X if: $Y = \varphi(X)$ for some φ s.t.

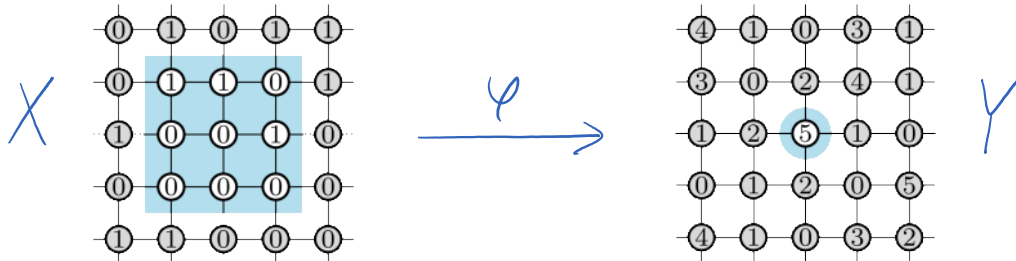
(*) φ is measurable

(*) φ is equivariant

[commutes with translations]
 $\varphi \circ T = T \circ \varphi \quad \forall T \text{ trans.}$

- φ is **finitary** if: Y_0 can be determined from the values of X on a finite box Λ_R (which may depend on the input X)

[i.e. for ν -a.e. x there exists $R < \infty$ s.t. $\varphi(x)_0$ only depends on $(x_u)_{u \in \Lambda_R}$]



- **Example:** X IID Bernoulli($1/2$)

- $Y_u := \min \{ \|u - v\|_\infty : X_u = X_v, u \neq v \}$

finitary factor

- $Y_u := \max \{ k \geq 0 : X_u = X_v \text{ for } \|u - v\|_\infty = k \}$

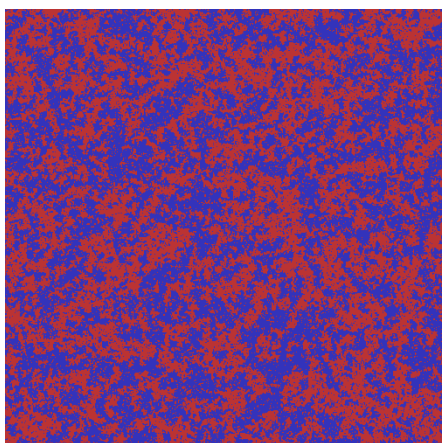
factor, non-finitary

- **Question:** Given Y , can Y be represented as

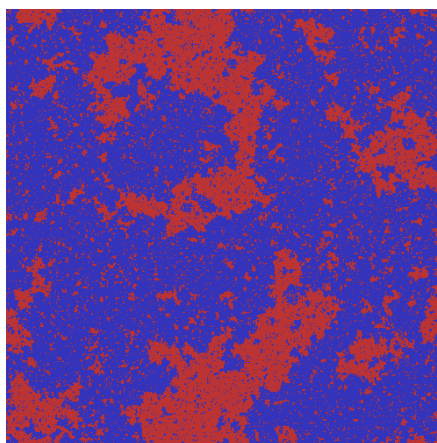
QUESTION: ... , can ... representation ...

- a factor of IID? (Bernoulli)
- a finitary factor of IID? (FFIID)
(and if so, is there an efficient coding?)

Ferromagnetic Ising model on \mathbb{Z}^d : ($d \geq 2$)



$\beta < \beta_c$



$\beta = \beta_c$



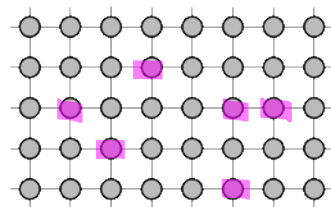
$\beta > \beta_c$

- μ^+ is a factor of IID for all β . (Ornstein, Weiss ~1977)
 \Rightarrow factor of IID cannot "detect" the phase transition.
- μ^+ is FFIID $\Leftrightarrow \mu^+ = \mu^-$ (van den Berg, Steif 1999)
[$\Leftrightarrow \beta \leq \beta_c$]

What is the algorithm that produces the finitary coding?

- Discrete-time Glauber dynamics

\Rightarrow Starting from any τ , we get $w_1^\tau, w_2^\tau, w_3^\tau, \dots$



- Coupling from the past (CFTP) (Propp-Wilson '96)

Instead of starting at time 0 and running to time n, start at time -n and run to time 0.

$$w_n^\tau = w_{0 \rightarrow n}^\tau \Rightarrow \hat{w}_n^\tau = w_{-n \rightarrow 0}^\tau$$

- Monotonicity

$$\Rightarrow \hat{w}_1^+ \geq \hat{w}_2^+ \geq \hat{w}_3^+ \geq \dots$$

$$\Rightarrow \hat{w}^+ := \lim_{n \rightarrow \infty} \hat{w}_n^+ \text{ exists almost surely}$$

- $\hat{w}^+ \sim \mu^+$, $\hat{w}^- \sim \mu^-$, $\hat{w}^- \leq \hat{w}^+$ a.s.

- $\mu^+ = \mu^- \Rightarrow \hat{w}^- = \hat{w}^+$ a.s.

$$\Rightarrow T := \min \{ |n| : \hat{w}_n^-(0) = \hat{w}_n^+(0) \}$$

is a.s. finite.

van den Berg and Steif

- $\beta = \beta_c$: any finitary coding has infinite expected coding volume: $\mathbb{E} R^d = \infty$, (+ Peres)
- $\beta < \beta_c$: can use finite-valued IID OR can get exponential tails for coding radius.

[uses that T has exp. tails (Martinielli-Olivieri 94)]

each answers

- $\beta < \beta_c$: can use finite-valued IID AND

each answers a question of van den Berg and Steif

- $\beta < \beta_c$: can use **finite-valued** IID **AND** (S. 2018)
can get **stretched-exp. tails** for coding radius.
 - $\beta = \beta_c$: can use finite-valued IID. (Meyerovitch; S. 2020+)
- [More generally, if Y is FF IID (and Y is finite-valued)
then can use finite-valued IID whose entropy is $h(Y) + \epsilon$.]
- $\beta > \beta_c$: "gradient" of m^+ is FF IID (Ray, S. 2019)
[implies that "energy" in a box has
volume-order large deviation estimate]

Other models: Potts, colorings, hard-core, random-cluster,

Meta-question (the van den Berg - Steif phenomenon):

When does FF IID \leftrightarrow phase transition?

Question (van den Berg and Steif 99): If a Markov random field is "unique", must it be a factor of IID?

Question (Steif): If a subshift of finite type has a unique measure of maximal entropy which is a factor of IID, must it also be FF IID?

taken from
"Open problems in symbolic dynamics"
Boyle 2008

Ferromagnetic Ising model

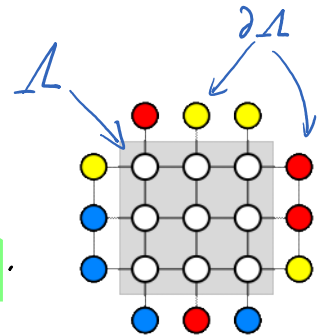


↙
 Nearest-neighbor interactions,
 Markov random field (MRF)
 [e.g. Potts, colorings]

↘
 Monotonicity, FKG
 [e.g. random-cluster model]

Markov random fields:

$X = (X_v)_{v \in \mathbb{Z}^d}$ is a MRF if given X_{Λ^c} ,
 the conditional law of X_{Λ} only depends on $X_{\partial\Lambda}$.



Suppose X is a translation-invariant MRF taking values in S .

Theorem (van den Berg, Steif 99, S. 2018)

Non-uniqueness $\Rightarrow X$ is not FFID
 +
 mild technical assumption [e.g. finite energy]

High noise:

$$\gamma := \sum_{s \in S} \min_{\tau \in S} \mathbb{P}(X_v = s \mid X_{2v} = \tau)$$

[called the multigamma admissibility]

Theorem (Häggström, Steif 2000)

$\gamma > 1 - \frac{1}{2d} \Rightarrow X$ is FF IID with exp. tails

stronger than Dobrushin uniqueness assumption

[e.g. the unique Potts measure when $|\beta| < \beta_0(q, d) \approx \frac{c(q)}{d^2}$]

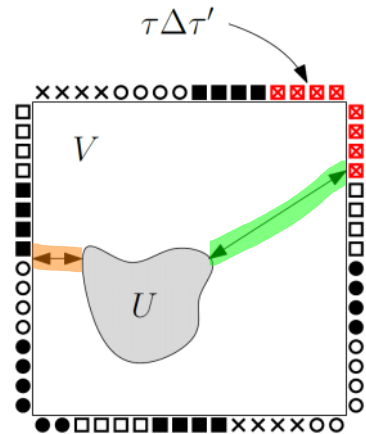
- Proper q -colorings : theorem does not apply since $\gamma=0$, but same result holds when $q \geq 4d(d+1)$. (Huber 99, S. 2018)

- Weak and strong spatial mixing

WSM : $\|p_{V,u}^\tau - p_{V,u}^{\tau'}\|_{TV} \leq C|u| e^{-c \text{dist}(u, \partial V)}$

SSM : $\|p_{V,u}^\tau - p_{V,u}^{\tau'}\|_{TV} \leq C|u| e^{-c \text{dist}(u, \tau \Delta \tau')}$

[Dobrushin uniqueness \Rightarrow SSM]



- Theorem (S. 2018)

◦ WSM $\Rightarrow X$ is FF IID with power-law tails

◦ SSM $\Rightarrow X$ is FF IID with exp. tails

[e.g. ferromagnetic Potts for all $\beta < \beta_c$]
 [e.g. proper q -colorings when $q \geq 3.53d$]

◦ WSM can be relaxed to $\exists \delta \in (0,1) \forall n \gamma_{\Lambda_n, \Lambda_n} \geq c > 0$.

where $\gamma_{V,U} = \sum_{w \in S^U} \min_{\tau \in S^{2V}} \mathbb{P}(X_u = w \mid X_{2V} = \tau)$

[e.g. Potts $d=2, \beta = \beta_c$ is FFID $\Leftrightarrow q \leq 4$]

THANK YOU !!!